

Functional programming learning path

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The experience of teaching

Traditionally, the A. P. Ershov Institute of Informatics Systems of the Siberian Branch of the Russian Academy of Sciences, in cooperation with Novosibirsk State University, conducts research and development of methods for teaching programming.

- **1990 (the beginning**) a **continuously developing** course "Functional Programming" on the basis of Mechanics and Mathematics Faculty of Novosibirsk State University.
- **1990 (the late**) courses "Programming Paradigms" (the sections on **functional programming**) for the Faculty of Mechanics and Mathematics, the Higher College of Informatics and the Information Technologies Faculty of Novosibirsk State University.

Distance courses

- 2004 Functional Programming Fundamentals https://intuit.ru/studies/courses/29/29/info
- 2006 Introduction to Lisp Programming

https://intuit.ru/studies/courses/1026/158/info.

- Programming paradigms

https://intuit.ru/studies/courses/1109/204/info

2019 - A new course "Programming Paradigm"

based on the Moodle system.

https://www.iis.nsk.su/files/book/file/FIT-Gor-PP3.pdf

Functional programming learning path

- laboratory practice on semantic principles,
- understanding of **pragmatics principles**, which frees the programmer from solving problems that are not fundamental for the nature of tasks,
- familiarization with a number of functional programming languages,
- the **principles formulation** of the functional programming paradigm that **distinguish** it from other paradigms.

Skills, understanding, outlook

Relationships between of principles,

consequences,

applications,

and practical trade-offs in functional programming.

The important idea: **the correctness of programmed decisions is more important** than the effectiveness of resulting programs is sharply emphasized.



Functional programming is one of the first paradigms

aimed not so much at obtaining an efficient implementation of pre-created and well-studied algorithms, but at a complete solution of **new and research problems.**

For functional programming, **priority is given to the organization of calculations and data structures**, and issues of memory processing and process control are pushed to the periphery of attention.

Semantic principles & consequence

Universality - to define debug-friendly generic functions, always to produce a result

- **Self-application** achieve clear, understandable and concise definitions through the use of recursion
- **Equal rights of parameters -** choose independent parameters with a clear representation of the boundary between bound and free variables
- **Meta-programming** to actively use and edit debugged templates
- **Verification** shows the role of proofs and axiomatization in ensuring the correctness of programs
- **The autonomy of developed modules** shows the possibilities of multidimensional combinatorics

Pragmatic principles & consequence

The flexibility of restrictions - can be reproduced as the processing of vectors.

Data immutability - can be shown on the mechanisms of continuous software development systems.

The rigorous results - show its advantages when familiarized with the definition of abstract machines.

Assumption of process continuity - the description of processes whose boundaries are difficult to predict.
 Reversibility of actions - debugging programs

dynamically, like UNDO in editing systems.

Unary functions - for turning program fragments into uniform flows of actions

Step 2

- The framework of functional programming allows one to consider **independent streams** of represented data and, if necessary, reorganize the stream space.
- With the transition to reusable programs and parallel computing, **the performance of program becomes more important** than their formal correctness.
- Production functional programming typically includes **practical trade-off mechanisms** that look like special functions in the programming language.

For parallel computing

Lazy evaluation - allows to avoid the cost of executing for too rare situations.

- **Identity of recalculations -** for debugging programs and measuring their performance.
- The iteration space in the absence of relationships between iterations (supported in the **Sisal** language).
- Automatic parallelization, load balancing, multi-pin nodes can be useful for solving problems of organizing parallel computing.

Most often automatic parallelization is supported.

Support for load balancing and multi-pin nodes is not common.

Productivity increase & Outside world

Data type control - analysis of data types.
Loop diagrams - familiar computational control scheme.
Memoization - radically can be reducing the complexity of calculations.

Programmable prediction - to think about the actual boundaries on available resources, including time.
Pseudo-functions - access to external devices for I/O.
Data recovery - correct alternative to data immutability, aimed in case of performance problems to transition debugged programs to more effective means without violating the correctness of their processing.

- Let us consider the following task: **the sum of cubes of nonnegative elements whose indices are even**.
- Let us note that the *map*, *reduce*, and *select* functions allows simplifying solutions of many similar problems and allow parallel execution of such solutions.
- A function-parameter of *map* and *select* should depend on element index to solve the considering problem.
- But function-parameter of *map* and *select* depends only on sequence element.
- Consequently, an important challenge is the generalization of the *map* and *select* functions.

Students consider this task and may suggest different solutions.

The solution of the problem is the iteration space from the **Sisal** programming language.

Students can execute Sisal programs using cloud parallel programming system (CPPS).

An iteration space can be created using Sisal triplets or the

Cartesian product of triplets.

A triplet defines an arithmetic progression.

The following construct is a triple construction expression:

lower_bound..upper_bound..step

Examples of triplets are the following constructs:

- •0..n.1 is a sequence of natural numbers up to *n*;
- •0..n-1..2 is a sequence of indices of even elements of an array of length *n*.

The following construct defines a Sisal **loop controlled by a range**: for var in triplet do returns reduction expr

end for

- •*var* is a variable;
- •*triplet* is a triplet;
- •*expr* is a reducible expression (it may depend on *var*);
- •*reduction* is a reduction (for example, *sum of*, *array of*).
- The **semantics** of this loop may be described as follows:
- Each iteration corresponds to each triplet element (*var* corresponds to the value of the triplet element).
- The value of a single iteration is the value of the reducible expression. 14
- The value of a loop is the value of a reduction.

The solution of the considering problem is the following function: function sum elements even indices(a: array of integers n: integer returns integer) for i in 0..n-1..2 do returns sum of $if(a[i] \ge 0)$ a[i]*a[i]*a[i] else 0 end if end for end function This task demonstrates comparation of *map-reduce-select* approach and Sisal loops to students.

Students can apply deductive verification approach to Sisal programs using the C-lighVer system

- We suggest students to use symbolic method of verification of definite iterations for Sisal program deductive verification.Symbolic method of verification of definite iterations is applied to special kinds of loops (definite iterations).
- Body of a definite iteration is executed once for each element of data sequence.

Symbolic method of verification of definite iterations allows defining inference rules for these loops without invariants.Symbolic replacement of definite iterations by recursive

functions (function \$rep\$) is the base of this method. We use ACL2 as theorem prover in the C-lightVer system.

Let us consider the precondition of considering function written in the language of ACL2 system:

(and (integerp n) (< 0 n) (equal (length a) n) (integer-listp a)) The postcondition is (= result (reduce-sum-even-indices n a)). Replacement of *result* term in postcondition by application of

rep results in the following verification condition:

(implies

(and (integerp n) (< 0 n) (equal (length a) n) (integer-listp a)) (equal

(rep (triplet 0 (- n 1) 2)) a)

(reduce-sum-even-indices n a)))

ACL2 successfully proved it by induction on *n*.

This verification task demonstrates correspondence between iterations and recursion to students. 17

Step 3

- With the **rapidly expanding space** of programming languages and the development of a **multitude of programming paradigms**, it is necessary to familiarize students with **the trends in functional programming**, its prospects, its place among other paradigms.
- Mostly these are historically significant and popular languages, including

Pure Lisp, Clisp, Haskell, F#, Clojure, and others.

A description of the actual functional programming techniques with an overview of functional programming languages and the relationship of functional programming with other paradigms.

Concluding Remarks

The focus of functional programming on solving many problems of organizing parallel computing.

Some questions have not yet received a practical answer.

Strictly speaking, **functional programming can play the role** of not only a training studio for the growth of professional programmer qualifications, but also a **design department** for production programming.

Any questions?

We thank you for your attention!

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- Gorodnyaya, L. Strategic Paradigms of Programming, Which Was Initiated and Supported by Academician Andrey Petrovich Ershov /Selected Papers 2020 5th International Conference on the History of Computers and Informatics in the Soviet Union, Russian Federation and in the Countries of Mutual Economic Assistance Council, SoRuCom 2020, 2020, p. 1–11
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Distance courses (in Russian)

2004 - Functional Programming Fundamentals

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- 2015 The educational and methodical **teaching aid** course "Programming Paradigm" https://www.iis.nsk.su/files/book/file/FIT-Gor-PP3.pdf
- 2019 The **new course** "Programming Paradigms" for FIT NSU based on the Moodle system https://el.nsu.ru/course/view.php?id=1200 (available at NSU)