

How fuzzing can help improve compiler quality

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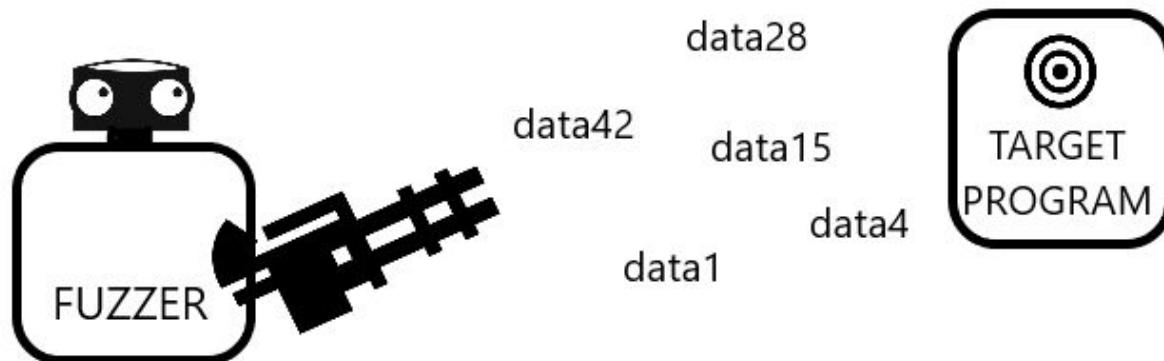


Problem

- All programs contains bugs
- Compilers are not an exclusion
 - Clang, GCC, javac, etc. bug trackers contains tens of thousands issues
- Manual testing is not enough
- Is there any way to improve compilers quality?

Fuzzing

“Program that generates a stream of random characters to be consumed by a target program”, Barton Miller et al., 1988



Fuzzers diversity

Black-box:

- Working with I/O of program

Grey-box:

- Gets partial information from the target program

White-box:

- Dynamic symbolic execution + coverage-maximizing heuristic

Disadvantages:

- Completeness

Disadvantages:

- Trying to find a balance between black-box and white-box fuzzing

Disadvantages:

- Implementation complexity
- All problems of symbolic execution

Black-box fuzzing example

```
fun getIthEl(a: Int, i: Int): Int {  
    val list = listOf(1, 2, 3)  
    return if (a == 1) {  
        if (i > 0) {  
            list[i]  
        } else -1  
    } else -1  
}
```



```
getIthEl(-1, 0)  
getIthEl(1, 2)  
getIthEl(-100, 21545135)  
getIthEl(125123590, 5123553)
```

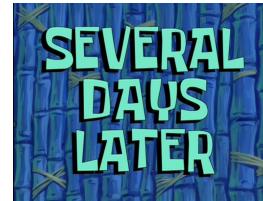


getIthEl(1, 123)

Grey-box fuzzing example

```
fun getIthEl(a: Int, i: Int): Int {  
    val list = listOf(1, 2, 3)  
    return if (a == 1) {  
        if (i > 0) {  
            list[i]  
        } else -1  
    } else -1  
}
```

getIthEl(-1, 0)
getIthEl(-1, -2142135)
getIthEl(-125, 352)
getIthEl(1, 0)



getIthEl(1, 1)
getIthEl(1, 123)

White-box fuzzing example

```
fun getIthEl(a: Int, i: Int): Int {  
    val list = listOf(1, 2, 3)  
    return if (a == 1) {  
        if (i > 0) {  
            list[i]  
        } else -1  
    } else -1  
}  
  
@S %0 = new java/lang/Integer[3]  
@S %1 = java/lang/Integer.class.valueOf(1)  
@S *(%0[0]) = %1  
@S %2 = java/lang/Integer.class.valueOf(2)  
@S *(%0[1]) = %2  
@S %3 = java/lang/Integer.class.valueOf(3)  
@S *(%0[2]) = %3  
@S %4 = kotlin/collections/CollectionsKt.class.listOf(%0)  
@S %5 = arg$0 != 1  
) -> (BEGIN  
<OR> (  
    @P %5 = false  
    @S %6 = arg$1 <= 0  
) -> (  
    @P %6 = false  
    @S %7 = %4.get(arg$1)  
    @S %8 = (%7 as java/lang/Number)  
    @S %9 = %8.intValue()  
    @S %10 = %9  
, <OR> (  
    @P %5 = false  
    @S %6 = arg$1 <= 0  
) -> (  
    @P %6 = true  
    @S %10 = -1  
, <OR> (  
    @P %5 = true  
    @S %10 = -1  
) END) -> (  
    @S <retval> = %10
```

```
[DEBUG] - Args: [2, 123]  
[DEBUG] - Collected trace: (  
    @S %0.1 = new java/lang/Integer[3]  
    @S %1.2 = java/lang/Integer.class.valueOf(1)  
    @S *(%0.1[0]) = %1.2  
    @S %2.3 = java/lang/Integer.class.valueOf(2)  
    @S *(%0.1[1]) = %2.3  
    @S %3.4 = java/lang/Integer.class.valueOf(3)  
    @S *(%0.1[2]) = %3.4  
    @S %4.5 = kotlin/collections/CollectionsKt.class.listOf(%0.1)  
    @S %5.6 = arg$0 != 1  
    @P %5.6 = true  
    @S %10.7 = -1  
    @S true = true  
)
```

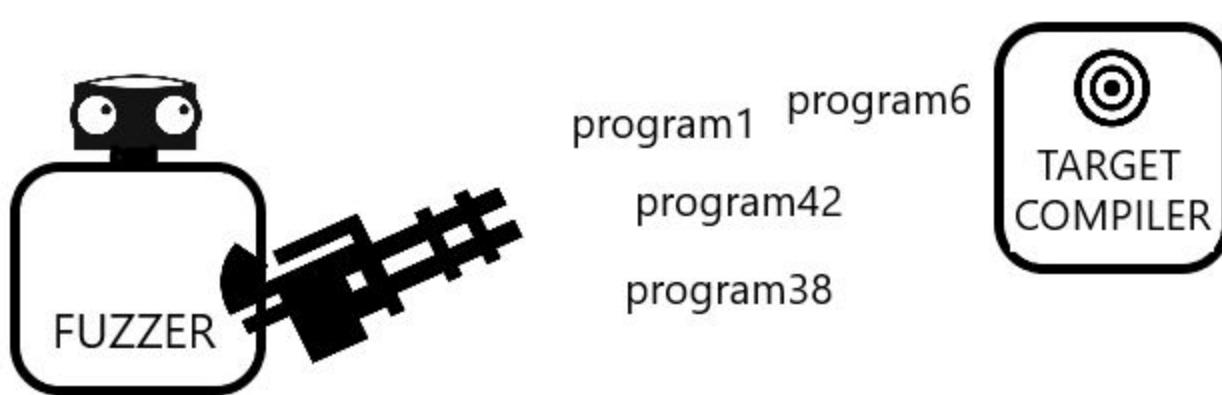
White-box fuzzing example

```
[DEBUG] - Args: [2, 3]
[DEBUG] - Collected trace: (
    @S %0.1 = new java/lang/Integer[3]
    @S %1.2 = java/lang/Integer.class.valueOf(1)
    @S *(%0.1[0]) = %1.2
    @S %2.3 = java/lang/Integer.class.valueOf(2)
    @S *%0.1[1]) = %2.3
    @S %3.4 = java/lang/Integer.class.valueOf(3)
    @S *(%0.1[2]) = %3.4
    @S %4.5 = kotlin/collections/CollectionsKt.class.listOf(%0.1)
    @S %5.6 = arg$0 != 1
    @P %5.6 = true
    @S %10.7 = -1
    @S true = true
)
```



Compiler fuzzing

“Program that generates a random programs to be consumed by a compiler”



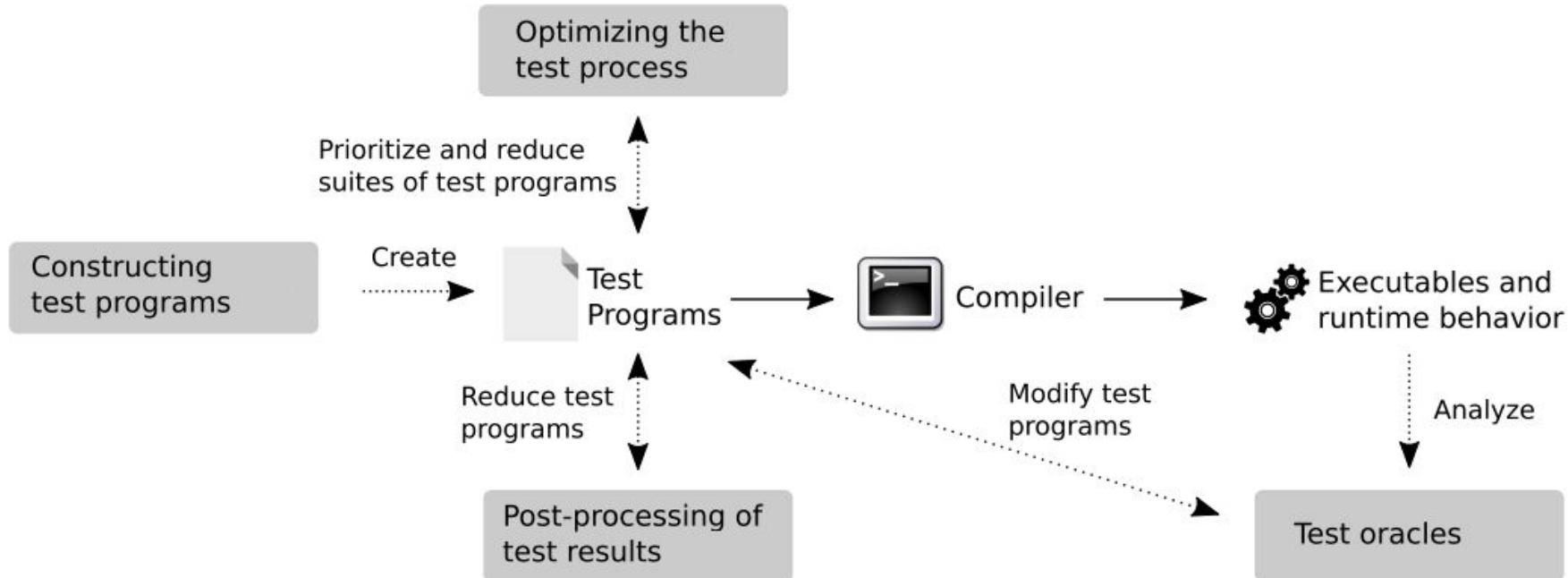
Compiler fuzzing

- Separate research area
- Started to develop in 1972
- Which language compilers are tested the most?
 - Javascript
 - C/C++
- Biggest complexity - input program generation

Fuzzers diversity

Language	Approach
C/C++	Eide and Regehr [47], Yang et al. [123], Nagai et al. [87], Nagai et al. [88], Lindig [78, 79], Groce et al. [51], Le et al. [72], Le et al. [73], Morisset et al. [85], Zhang et al. [131], Sun et al. [111], Amodio et al. [4], Alipour et al. [3]
JavaScript	Holler et al. [58], Groce et al. [49], Patra and Pradel [93], Bastani et al. [8]
PL/I	Hanford [52]
Ada	Duncan and Hutchison [45], Austin et al. [6], Mandl [80]
Fortran	Callahan et al. [17], Dongarra et al. [44], Burgess and Saidi [16]
Java	Sirer et al. [108], Boujarwah et al. [14], Yoshikawa [125], Chen et al. [31]
Algol	Houssais [61]
APL	Ching and Katz [32]
Arden	Wolf et al. [120]
Haskell	Palka et al. [92]
Lusture	Garoche et al. [48]
mpC	Kalinov et al. [64, 65]
OCaml	Midgaard et al. [83]
Pascal	Burgess [34], Bazzichi and Spadafora [9]
PLZ/SYS	Bazzichi and Spadafora [9]
Python	Bastani et al. [8]
Ruby	Bastani et al. [8]
Rust	Dewey et al. [41]
Scala	Zhang et al. [131]
GLSL	Donaldson et al. [42]

Compiler fuzzing problems



Methods of input programs construction

- Generational approach
 - Grammar-directed
 - Grammar-aided
 - Others
- Mutational approach
 - Semantics-preserving
 - Non-semantics-preserving

Grammar-directed approach



Purdom P. A sentence generator for testing parsers //BIT Numerical Mathematics.
– 1972. – T. 12. – №. 3. – C. 366-375.

Example of modern grammar-directed approach (LaLa specification*)

```
use SymbTab;

class Program {
    prog("${stmts : StmtList}\n") {
        stmts.symbols = (SymbTab:empty);
    }
}
@list
class StmtList {
    inh symbols : SymbTab;
    one("${s: Stmt}") {
        s.s_before = this.symbols;
    }
    @weight(3)
    mult("${s: Stmt}\n${r : StmtList}") {
        s.s_before = this.symbols;
        r.symbols = s.s_after;
    }
}

class Stmt {
    inh s_before : SymbTab;
    syn s_after : SymbTab;
    assign("${i : Identifier} := ${e : Expr};") {
        e.symbols = this.s_before;
        this.s_after = (SymbTab:put this.s_before i.str);
    }
}

class Expr {
    inh symbols : SymbTab;
    num("${val : Number}") {}
    use_var("${var : UseVariable}") {
        var.symbols = this.symbols;
    }
    binop("${l : Expr} ${op : Op} (${r : Expr})") {
        l.symbols = this.symbols;
        r.symbols = this.symbols;
    }
}

class UseVariable {
    inh symbols : SymbTab;
    grd defined;
    var("${var : Identifier}") {
        this.defined =
            (SymbTab:contains this.symbols var.str);
    }
}

# alternatively:
class UseVariableGen {
    inh symbols : SymbTab;
    var(SymbTab:defs this.symbols) : String {}
}

class Identifier("[a-z]{1,3}");
class Number("0|[1-9][0-9]{0,2}");
class Op("+|-|*|/");
```

* P. Kreutzer, S. Kraus, and M. Philippse, “Language-agnostic generation of compilable test programs,” in 2020 IEEE 13th International Conference on Software Testing, Validation and Verification (ICST). IEEE, 2020, pp. 39–50.

Grammar-aided approaches

- CSmith
 - Generator of programs on a subset of the C language
 - Uses grammar + a lot of analyses to ensure test validity

Problem	Code-Generation-Time Solution	Code-Execution-Time Solution
use without initialization	explicit initializers, avoid jumping over initializers	—
qualifier mismatch	static analysis	—
infinite recursion	disallow recursion	—
signed integer overflow	bounded loop vars	safe math wrappers
OOB array access	bounded loop vars	force index in bounds
unspecified eval. order of function arguments	effect analysis	—
R/W and W/W conflicts	effect analysis	—
betw. sequence points		
access to out-of-scope stack variable	pointer analysis	—
null pointer dereference	pointer analysis	null pointer checks

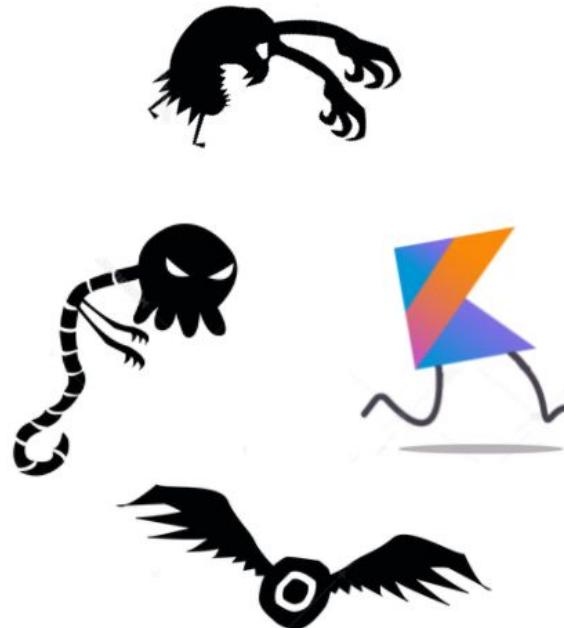
Table 1. Summary of Csmith’s strategies for avoiding undefined and unspecified behaviors. When both a code-generation-time and code-execution-time solution are listed, Csmith uses both.

Other approaches

- A lot of approaches, implements different “local” ideas:
 - Testing certain features in compilers by writing a generator of the desired language constructions
 - Machine learning based
 - Based on the derivation of the rules for generating programs from projects in the target language
 - ...

Mutation based approaches

- We are mutating existing part of program
 - Semantic-preserving
 - Non-Semantic-preserving



Semantics-preserving

- Metamorphic testing
 - We accidentally mutate a part of the program that is not covered at runtime
 - Insert “dead code”
 - Replacing expressions with equivalent ones

```
fun getIthEl(a: Int, i: Int): Int {  
    val list = listOf(1, 2, 3)  
    return if (a == 1) {  
        if (i > 0) {  
            list[i]  
        } else -1  
    } else -1  
}
```



```
fun getIthEl(a: Int, i: Int): Int {  
    val list = listOf(1, 2, 3)  
    return if (a == 1) {  
        if (i > 0) {  
            if (i == 0) -1  
            else list[i]  
        } else -1  
    } else -1  
}
```

Non-semantics preserving

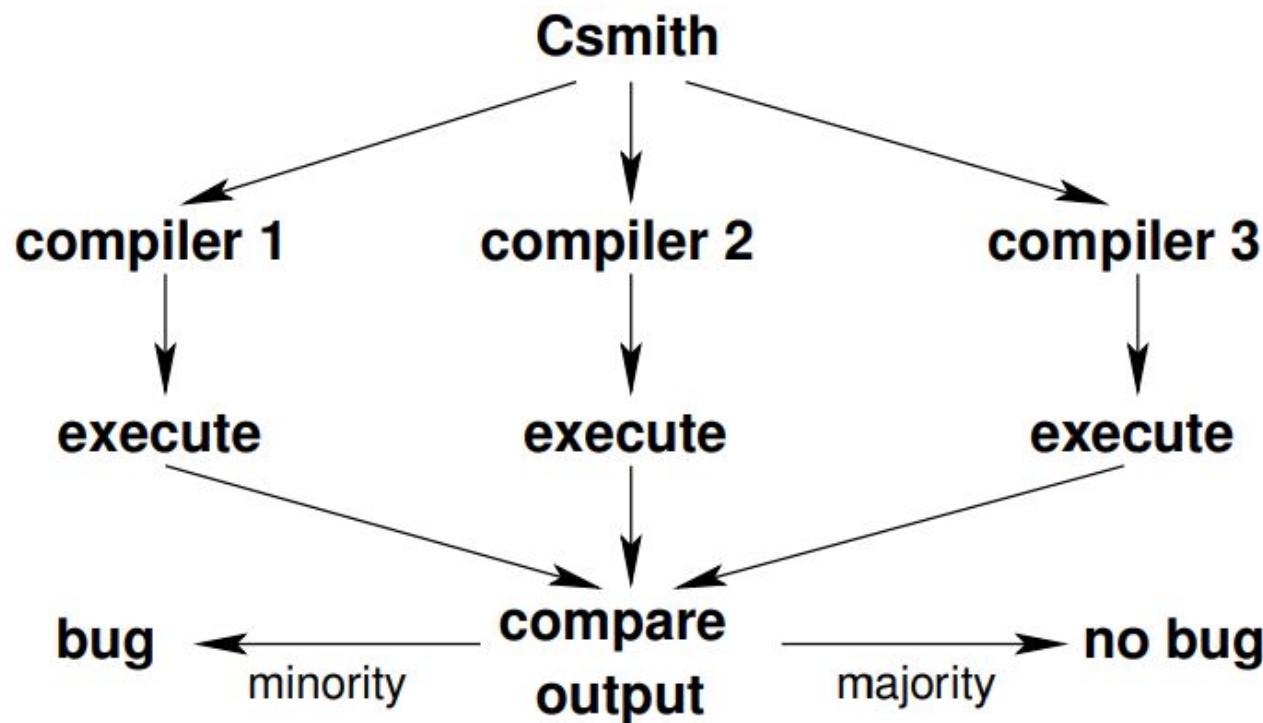
- Different mutations over different kinds of code representation
- If resulting test is incorrect, roll back

```
fun getIthEl(a: Int, i: Int): Int {  
    val list = listOf(1, 2, 3)  
    return if (a == 1) {  
        if (i > 0) {  
            list[i]  
        } else -1  
    } else -1  
}
```



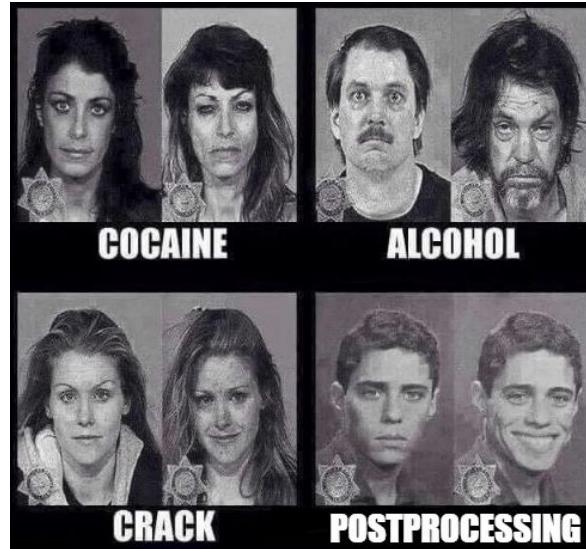
```
fun getIthEl(a: Int, i: Int): Int {  
    val list = listOf(1, 2, "3")  
    return if (a > 1) {  
        if (i == 0) {  
            list[i]  
        } else -1  
    } else -1  
}
```

Test oracles problem. How to detect miscompilation?



Postprocessing

- Reduction
- Duplicates filtration
- Bug isolation



Reduction. Problem

```
open class A1(y: String) {
    val x = "[A1.x,$y]"
}
open class A2(y: String) {
    val x = "[A2.x,$y]"
    inner open class B1 : A1 {
        constructor(p: String) : super("[B1.param,$p]")
        fun foo() = x + ";" + this@A2.x + ";"
    }
    fun bar(): String {
        return with(A2("#bar")) {
            class C : B1("bar") {}
            C().foo()
        }
    }
    fun foo() = A2("#baz").baz()
    fun A2.baz(): String {
        class C : B1("baz") {}
        return C().foo()
    }
}
fun box(): String {
    val r3 = A2("f").bar()
    if (r3 != "[A1.x,[B1.param,bar]];[A2.x,#bar];") return "$r3:"
    val r4 = A2("gg").foo()
    if (r4 != "[A1.x,[B1.param,baz]];[A2.x,#baz];") return "$r:"
    return "OK"
}
```

Reduction. Problem

```
open class A1

class A2 {
    open inner class B1 : A1()

    fun A2.baz() {
        class C : B1()
    }
}
```

Reduction. Existing methods

- Language-agnostic methods
 - Delta-debugging
 - Slicing
 - Hierarchical delta-debugging
 - ...
- Language-specific methods
 - Transformations over different types of code representation
- Different combinations of approaches
 - CReduce

Duplicates filtration. For what?

```
/*1*/ fun loop() { {}!!!! > 1 }
```

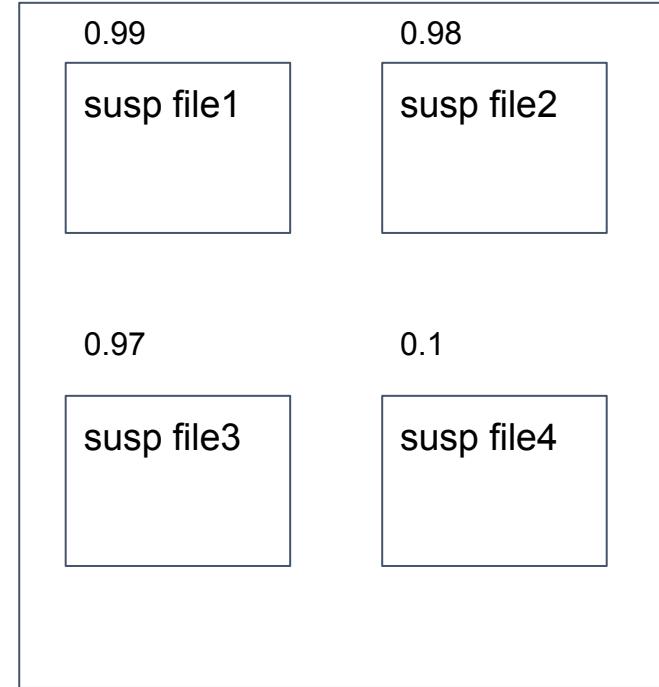
```
/*2*/ var a = {}!!!!
```

```
/*3*/ fun box() { -{}!!!! }
```

```
/*4*/ fun box() { ({}!!!! == 1).not() }
```

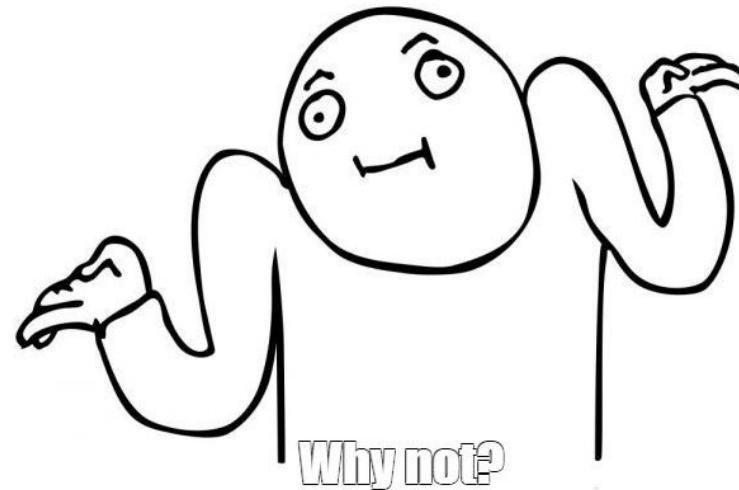
```
// ...
```

Bug isolation



BBF(Backend bug finder). Origins

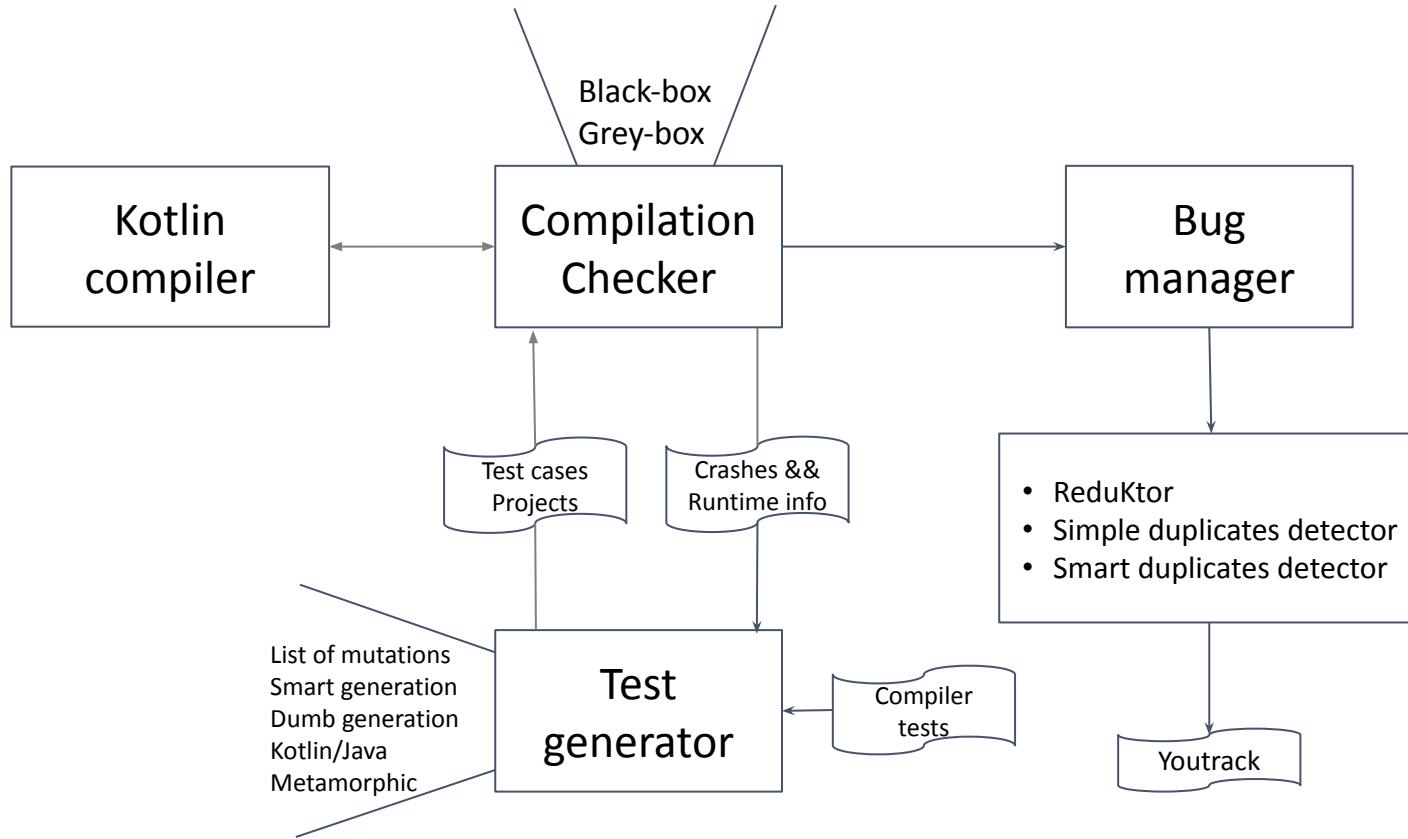
- 2017 - simple fuzzer
- 2018 - reducer
- 2018 -



What types of bugs can BBF find?

- Crashes
- Miscompilations
- Performance issues
- Bugs in plugin

BBF scheme



Why not generative approach

- The more different “features” in the language, the more difficult it is to generate it:
 - Smart cast-s
 - Tailrec
 - Null safety
 - Reflection
 - Interacting with libraries
 - ...
- Either we generate only a subset of the language, or we spend a lot of time on each feature

Test generator

- 30+ mutators of programs
 - Adding an argument to a function
 - Changing modifiers
 - Replacing a node in AST with a node of the same type
- Type-centric enumeration
- Metamorphic testing
- Random data structures generator

Skeletal program enumeration

```
var a: Int = 1
var b: Int = 1
```

```
while (b < 100) {
    val c = b
    b = a + b
    a = c
}
```

```
var [__]: Int = 1
var [__]: Int = 1
```

```
while ([__] < 100) {
    val [__] = [__]
    [__] = [__] + [__]
    [__] = [__]
}
```

(a) P

```
var a: Int = 1
var b: Int = 1

while (a < 100) {
    val c = a
    b = b + b
    b = a
}
```

(b) P'

```
var a: Int = 1
var c: Int = 1

while (a < 100) {
    val b = a
    c = a + c
    a = c
}
```

* Zhang Q., Sun C., Su Z. Skeletal program enumeration for rigorous compiler testing //Proceedings of the 38th ACM SIGPLAN Conference on Programming Language Design and Implementation. – 2017. – С. 347-361.

Type-centric fuzzing

```
class A {  
    val a: Int  
}  
  
fun f(a: Int) {  
    ...  
}  
  
var a: Int = 1  
var b: Int = 1  
  
while (b < 100) {  
    val c = b  
    b = a + b  
    a = c  
}
```

```
class A {  
    val a: Int  
}  
  
fun f(a: Int) {  
    ...  
}  
  
var a: Int = [Int]  
var b: Int = [Int]  
  
while ([Int] < [Int]) {  
    val c = [Int]  
    [Int] = [[Int] + [Int]]  
    [Int] = [Int]  
}
```

(a) P

```
class A {  
    val a: Int  
}  
  
fun f(a: Int) {  
    ...  
}  
  
var a: Int = -100  
var b: Int = a * f(a)  
  
while (a < b) {  
    val c = 12345  
    b = A(123).a  
    b = b  
}
```

(b) P'

```
class A {  
    val a: Int  
}  
  
fun f(a: Int) {  
    ...  
}  
  
var a: Int = A(7).a  
var b: Int = 3  
  
while (b < 100) {  
    val c = f(a * 10)  
    b = c * (-10) + a  
    a = b + c  
}
```

Type-centric fuzzing (generation phase)

INPUT: file with a seed program P
OUTPUT: list of generated expressions $c_0 \dots c_N$

```
1: function GENERATIONPHASE(file)
2:   res ← []
3:   for callee ∈ getCallables(file) do
4:     if callee is Class then
5:       instance ← genClassInstance(callee, file)
6:       for iCallee ∈ getCallables(instance) do
7:         res ← res + genCall(iCallee, file)
8:     else
9:       res ← res + genCall(callee, file)
10:    return res
11: end function
12:
13: function GENCLASSINSTANCE(class, file)
14:   typeParams ← genTypeParams(class, file)
15:   class ← parameterize(class, typeParams)
16:   if ¬hasOpenConstructor(class) then
17:     impl ← findImplementation(class)
18:     if impl ≠ null then
19:       impl ← adaptTypeParams(
20:         impl, class, typeParams)
21:       return genClassInstance(impl, file)
22:     else
23:       return null
24:   randomCtor ← getRandomConstructor(class)
25:   return genConstructorCall(randomCtor, file)
26: end function
```

Fig. 4: Generation phase algorithm

```
class A(val a: Int) {
  fun f(a: String): Int { ... }
}

interface B {
  val a: Int
}

class C(override val a: Int) : B

fun f(a: Int): Int

val a: Int = 1
```

(a) Seed program for generation phase

```
A(1) -> A
A(1).a -> Int
A(1).f("") -> Int
C(1) -> B
C(1) -> C
C(1).a -> Int
f(1) -> Int
a -> Int
```

(b) Generated expressions with their types

Fig. 3: Generation phase example

Type-centric fuzzing (mutation phase)

INPUT: seed program for mutation phase *seed*
INPUT: seed program from generation phase *gen*
INPUT: generated expressions *exprs*
OUTPUT: program will filled type placeholders

```

1: function MUTATIONPHASE
2:   anon ← anonymize(seed)
3:   anon ← merge(anon, gen)
4:   for ph ∈ getPlaceholders(anon) do
5:     e ← genPhExpr(ph, exprs)
6:     anon ← replacePhWithExpr(anon, ph, e)
7:   return anon
8: end function

9:
10: function GENPHEXPR(ph, exprs)
11:   type ← getType(ph)
12:   r ← []
13:   r ← r + genRandomValue(type)
14:   r ← r + genStdLib(type)
15:   for e ∈ exprs do
16:     eType ← getType(ph)
17:     if compatible(type, eType) then
18:       r ← r + e
19:   return random(r)
20: end function
```

```

val a: Int = 1

class A(val a: Int) {
  fun f(a: String): Int { ... }
}

fun f(a: Int): Int { ... }

// Generated expressions with their types:
// A(1) -> A
// A(1).a -> Int
// A(1).f("") -> Int
// f(1) -> Int
// a -> Int
```

(a) Seed program after generation phase

```

fun factorial(n: Int): Double {
  var result = 1.0
  for (i in 1..n) {
    result *= i
  }
  return result
}
```

(b) Seed program for mutation phase

```

fun factorial(n: Int): Double {
  var result = [Double]
  for (i in [IntRange]) {
    [Double] *= [Int]
  }
  return [Double]
}
```

(c) Typed skeleton

```

val a: Int = 1

class A(val a: Int) {
  fun f(a: String): Int { ... }
}
```

```

fun f(a: Int): Int { ... }
```

```

fun factorial(n: Int): Double {
  var result = f(1).toDouble()
  for (i in A(1).f("")..a) {
    result *= i
  }
  return result
}
```

(d) Typed skeleton filled with the generated expressions

Fig. 5: Mutation phase algorithm

Evaluation

	NO SEVERITY	MINOR	NORMAL	MAJOR
Frontend	2	0	0	0
Backend	0	1	5	4
Miscompilation	0	0	2	4

TABLE I: Posted bug severity from the Kotlin YouTrack

Results	M	G	EM	SPE	TCE	TCE + EM
Correct programs, %	11,9	0,05	10,7	24,2	63,4	13,4
Interesting bugs, %	25,0	0,0	20,0	100,0	100,0	16,6
Frontend crashes	3	212	4	0	0	5
Backend crashes	9	0	11	2	3	7
Miscompilations	0	0	0	0	3	0
Duplicates	49	77	38	1	4	14

TABLE II: Evaluation results

- (M) Mutation-based fuzzing;
- (G) Grammar based generation
- (EM) M + language-specific mutations
- (SPE) Skeletal program enumeration
- (TCE) Type-centric enumeration;
- (TCE + EM) TCE + language-specific mutations.

Black-box vs Grey-box target fuzzing

- Coverage is almost useless for standard testing
- Is the same also true for feature fuzzing?

JVM_IR indy-lambdas: initial implementation and tests

KT-44278 KT-26060 KT-42621

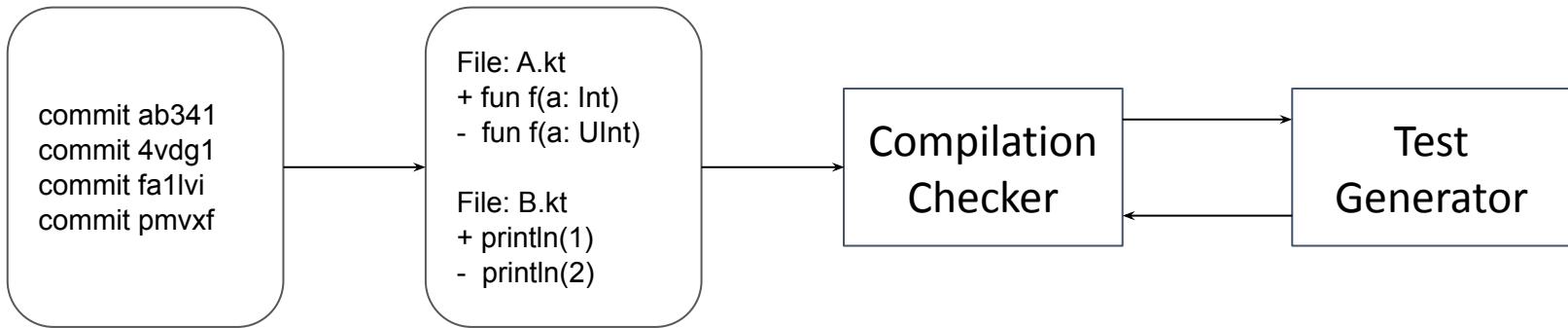
master build-1.5.20-dev-3945 ... build-1.5.20-dev-205

dnpetrov authored and TeamCityServer committed on Feb 3 1 parent 0bc386c commit d94912ed624d0347e4432ce69cc5465cad0ebe05

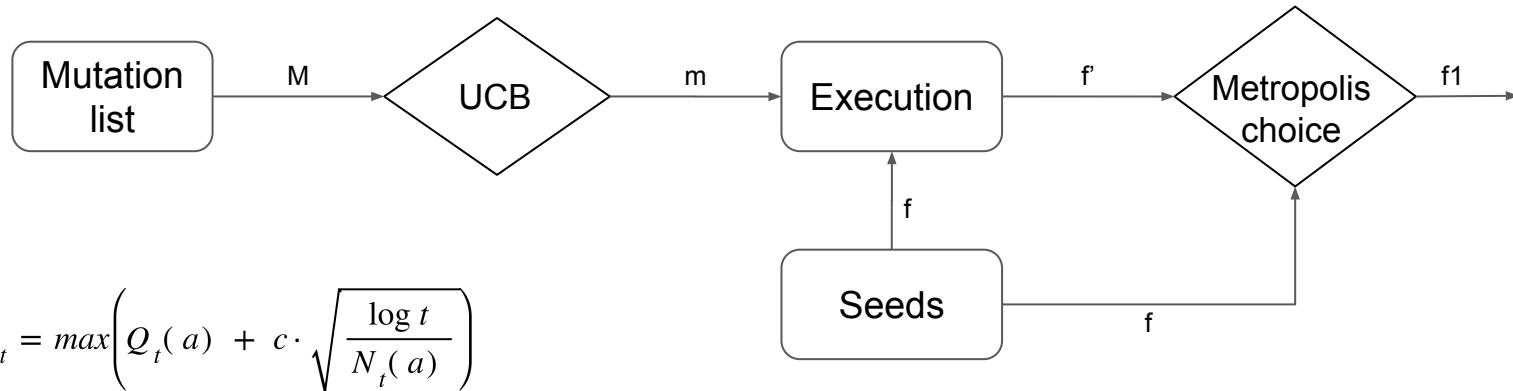
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Unified Split

Grey-box target fuzzing (main idea)



Grey-box target fuzzing (scheme)



$Q_t(a)$ – average points

t – iteration

$N_t(a)$ – number of "a" selections

$Q_t(a)$ – exploitation

$N_t(a)$ – exploration

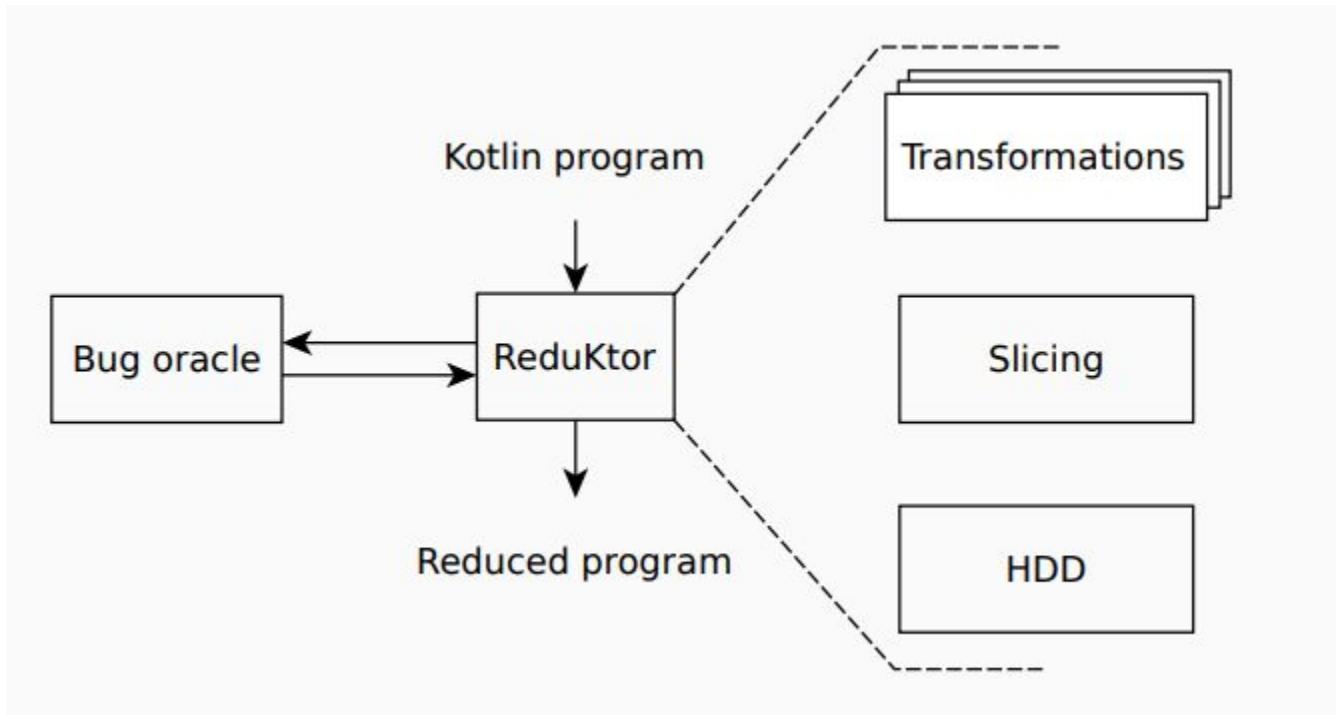
$$\text{Score}_{m_t} = \left(\text{coverage}(m_t) - \text{coverage}_{m_{t-1}} \right) \cdot K$$

$$A(f - >f_1) = \min(1, \exp(K \cdot (\text{cov}(f) - \text{cov}(f_1))))$$

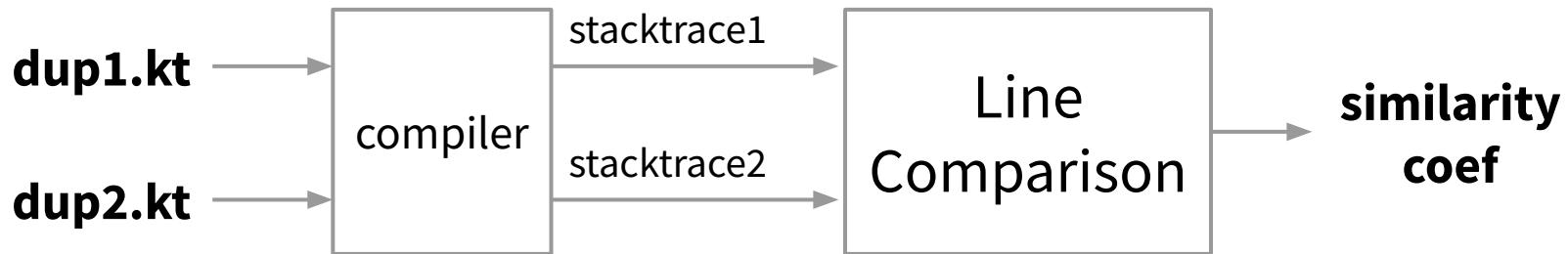
Postprocessors

- Reduction
- Duplicates filtration
- Bug isolation?

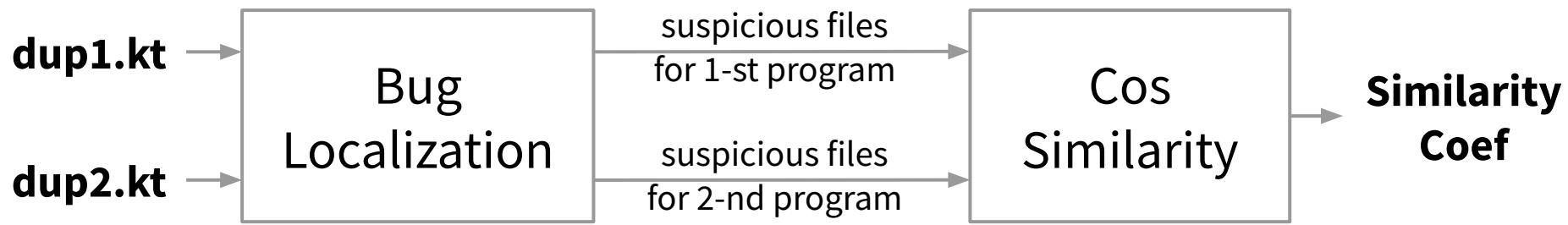
Scheme of ReduKtor



How duplicates filtration works?



Better duplicates filtration?



Results

- 213 bugs were posted found by BBF
 - 50 major
 - 39 normal
 - 9 minor
 - 88 not specified
- 83 fixed
 - 43 major
 - 20 normal
 - 5 minor
 - 15 not specified
- 34 open
- 6 in progress

How we helped to the Kotlin team?

- New backend ABI fuzzing
- Performance testing
- Fuzzing of new features
- Regressions search
- Kotlin plugin testing

Contacts

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